

Cycle tracking polynomial of a graph

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Abstract

We introduce a cycle tracking polynomial of a graph G . The cycle tracking polynomial of a graph G of order n is the polynomial $T(G, x)$ of G is defined as

$$T(G; x) = \sum_{i=\tau(G)}^n t(G, i)x^i$$

where $t(G, i)$ is the number of cycle tracking sets of G of size i , and $\tau(G)$ is the cycle tracking number of G . We obtain some properties of $T(G, x)$ and its coefficients. Also we compute this polynomial for some specific graphs.

AMS Subject Classification: 05C50, 05C69

Key Words and Phrases: cycle tracking set, cycle tracking number $\tau(G)$, cycle tracking polynomial

1 Introduction

The concept of cycle tracking set is introduced in [4] Let $G(V, E)$ be a graph. For $v \in V(G)$, the *cycle trace (simply trace)* of v is defined as the set of all vertices $u \in V$ such that u and v belong to same cycle of G and is denoted by $T_G(v)$. For $u, v \in V$, v is said to be *cycle traced (traced)* by u if $v \in T_G(u)$. A set S of vertices in a graph $G(V, E)$ is called a *cycle tracking set* if for every vertex $v \in V \setminus S$, there exists a vertex $u \in S$ such that $v \in T_G(u)$. A cycle tracking set is a *minimal cycle tracking set* if no proper subset $S' \subset S$ is a cycle tracking set. The *cycle tracking number*

$\tau(G)$ of a graph G is the minimum cardinality of a minimal cycle tracking set of G . The *upper cycle tracking number* $T(G)$ of a graph G is the maximum cardinality of minimal cycle tracking set of G . A cycle tracking set with minimum cardinality is called a τ – set of G .

A graph $G(V,E)$ is said to be *transitive tracking* graph if for every $u, v, w \in V(G)$, $w \in T_G(u)$ and $u \in T_G(v)$ implies $w \in T_G(v)$. A graph $G(V,E)$ is said to be *track connected* if for every pair of vertices $u, v \in V(G)$ there exist two internally disjoint paths connecting u and v . A vertex is said to be *trace free vertex* if it belongs to no cycle.

Through out this paper the letter G denotes a graph of order n .

2 Cycle tracking polynomial of a graph

Definition 2.1. Let G be a graph of order n . Let $T(G, i)$ be the family of all cycle tracking sets of a graph G with cardinality i and let $t(G, i) = |T(G, i)|$. Then the cycle tracking polynomial $T(G, x)$ of G is defined as

$$T(G; x) = \sum_{i=\tau(G)}^n t(G, i)x^i$$

where $\tau(G)$ is the cycle tracking number of G .

The path P_3 on three vertices has only one cycle tracking set with cardinality 3 ($\tau(G) = 3$) its tracking polynomial is then $T(P_3, x) = x^3$. In the case of the cycle C_n on $n(n \geq 0)$ vertices,

$$T(C_n, x) = \binom{n}{1}x^1 + \binom{n}{2}x^2 + \binom{n}{3}x^3 + \dots + \binom{n}{n}x^n = (1 + x)^n - 1.$$

Theorem 2.1. If a graph G consist of m components G_1, G_2, \dots, G_m then $T(G, x) = T(G_1, x)T(G_2, x)\dots T(G_m, x)$.

Proof. It is enough to prove the theorem for $n=2$.

For $k \geq \tau(G)$, a cycle tracking set of k vertices in G arises by choosing a cycle tracking set of j vertices in G_1 for some j such that $\tau(G) \leq j \leq |V(G)|$ and a cycle tracking set of $k - j$ vertices in G_2 . The number of ways of doing this over all $j = \tau(G_1), \dots, |V(G_1)|$ is exactly the coefficient of x^k in $T(G_1, x)T(G_2, x)$. So $T(G, x) = T(G_1, x)T(G_2, x)$. \square

Theorem 2.2. Let G be a graph of order n . Then

1. $t(G, n) = 1$
2. $t(G, i) = 0$ if and only if $i < \tau(G)$ or $i > n$
3. $T(G, x)$ has no constant term.
4. $T(G, x)$ is a strictly increasing function on $(0, \infty)$

- 5. for any subgraph H of G , $deg(T(G, x)) \geq deg(T(H, x))$
- 6. zero is a root of $T(G, x)$ with multiplicity $\tau(G)$
- 7. $\tau(G) = n$ if and only if $T(G, x) = x^n$

Theorem 2.3. [4] Let G be a graph of order n . Then $\tau(G) = n$ if and only if $G(V, E)$ is a forest.

Theorem 2.4. Let G be a graph of order n . Then $T(G, x) = x^n$ if and only if G is a forest.

Proof. $T(G, x) = x^n$ if and only if $V(G)$ is the only cycle tracking set for G . That is if and only if $\tau(G) = n$. That is if and only if G is a forest (by Theorem 2.3)

□

Theorem 2.5. Let G be a graph of order n . Then $T(G, x) = (1 + x)^n - 1$ if and only if G is track connected.

Proof. If G is track connected then $\tau(G) = 1$ and every vertex traces all vertices of G . So coefficient of x is n and $t(G, p) = \binom{n}{p}$. So $T(G, x) = \binom{n}{1}x^1 + \binom{n}{2}x^2 + \binom{n}{3}x^3 + \dots + \binom{n}{n}x^n = (1 + x)^n - 1$.
Conversely if $T(G, x) = (1 + x)^n - 1$, the coefficient of x is n . That is, every vertex traces all vertices of G . Hence G is track connected.

□

Definition 2.2. [4]

Let G be a graph with exactly one cut vertex. Let v be the cut vertex of G and G_1, G_2, \dots, G_k be the components of $G \setminus \{v\}$. If the order of G_i is greater than or equal to three and the graphs induced by $V(G_i) \cup \{v\}, i = 1, 2, \dots, k$ are track connected then G is called a *track connected floral graph*. And the graph induced by $V(G_i) \cup \{v\}, i = 1, 2, \dots, k$ are called petals of G .

Theorem 2.6. Let G be a track connected floral graph with k petals with each petal having m_1, m_2, \dots, m_k vertices respectively then

$$t(G, p) = \binom{m_1 + m_2 + \dots + m_k}{p-1} \text{ when } 1 \leq p \leq k-1 \text{ and}$$

$$t(G, p) = \binom{m_1 + m_2 + \dots + m_k}{p-1} + \sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \right. \right. \\ \left. \left. \left[\dots \left[\sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \dots \right] \right] \right]$$

when $k \leq p \leq n$ and

$$T(G, x) = \sum_{p=1}^{k-1} \binom{m_1 + m_2 + \dots + m_k}{p-1} x^p + \sum_{p=k}^n \left[\binom{m_1 + m_2 + \dots + m_k}{p-1} + \dots \right]$$

$$\sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \left[\dots \left[\sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \dots \right] \right] \right] x^p.$$

Proof. Case(1) $1 \leq p \leq k - 1$

Then any cycle tracking set S contains the central vertex. So the central vertex together with p-1 vertices constitute a cycle tracking set S and it can be chosen in $\binom{m_1+m_2+\dots+m_k}{p-1}$ ways.

Case(2) $k \leq p \leq n$

Here the central vertex together with p-1 vertices constitute a cycle tracking set S and a set of vertices S of cardinality p having at least one element from each leaf is also form a cycle tracking set and it can be chosen in

$$\binom{m_1 + m_2 + \dots + m_k}{p - 1} + \sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \left[\dots \left[\sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \right] \right] \right] \dots \Big] \text{ ways.}$$

□

Let G be a transitively tracked graph then its vertex set can be partitioned into V_1, V_2, \dots, V_k of cardinality m_1, m_2, \dots, m_k respectively such that the graph $\langle V_i \rangle$ induced by each V_i is maximal track connected subgraph of G. Then a set S of vertices which contains at least one element from each V_i form a cycle tracking set. So a cycle tracking set of G with cardinality p

can be chosen in $\sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \left[\dots \left[\sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \right] \right] \right] \dots \Big] \text{ ways.}$

So $t(G, p) = \sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \left[\dots \left[\sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \right] \right] \right] \dots \Big]$,

$k \leq p \leq n$.

The above discussion may be summarized as follows.

Theorem 2.7. *Let G is transitively traced and let $V(G)$ is partitioned into V_1, V_2, \dots, V_k of cardinality m_1, m_2, \dots, m_k respectively such that the graph $\langle V_i \rangle$ induced by each V_i is maximal track connected subgraph of G. Then*

$$T(G, x) = \sum_{p=k}^n \left[\sum_{i_1=1}^{p-k+1} \binom{m_1}{i_1} \left[\sum_{i_2=1}^{p-k-i_1+2} \binom{m_2}{i_2} \left[\dots \sum_{i_j=1}^{p-k-i_1-i_2-\dots+j} \binom{m_j}{i_j} \left[\dots \sum_{i_{k-1}=1}^{p-k-i_1-i_2-\dots-i_{k-1}+k-1} \binom{m_{k-1}}{i_{k-1}} \binom{m_k}{p-i_1-i_2-\dots-i_{k-1}} \right] \right] \right] \right] x^p.$$

Theorem 2.8. [4] Let G' be the graph formed by removing all cut edges of a graph G . Then a subset S of $v(G)$ is a cycle tracking set of G if and only if S is a cycle tracking set of G' .

Theorem 2.9. [4] Let G be a transitively tracked graph. Then the components of G obtained by deleting all cut edges of G are precisely the maximal track connected subgraph of G .

Theorem 2.10. Let G be transitively tracked and let $V(G)$ be partitioned into V_1, V_2, \dots, V_k such that the graph $\langle V_i \rangle$ induced by each V_i is maximal track connected subgraph of G . Then $T(G, x) = T(\langle V_1 \rangle, x)T(\langle V_2 \rangle, x)\dots T(\langle V_m \rangle, x)$.

Proof. Let G be transitively tracked and let $V(G)$ be partitioned into V_1, V_2, \dots, V_k such that the graph $\langle V_i \rangle$ induced by each V_i is maximal track connected subgraph of G . Let G' be the graph formed by removing all cut edges of G . Then by Theorem 2.8 $T(G, x) = T(G', x)$ and by Theorems 2.9 and 2.1 $T(G, x) = T(\langle V_1 \rangle, x)T(\langle V_2 \rangle, x)\dots T(\langle V_m \rangle, x)$ \square

Corollary 2.11. Let G be transitively tracked and let $V(G)$ be partitioned into V_1, V_2, \dots, V_k of cardinality m_1, m_2, \dots, m_k respectively such that the graph $\langle V_i \rangle$ induced by each V_i is maximal track connected subgraph of G . Then $T(G, x) = ((x + 1)^{m_1} - 1)((x + 1)^{m_2} - 1)\dots((x + 1)^{m_k} - 1)$

Theorem 2.12. For a graph G , $t(G, 1) = 1$ if and only if G is a track connected floral graph.

Proof. Let G be any graph with $t(G, 1) = 1$. Then there exists one and only one cycle tracking set S with $|S| = 1$. That is there exist a vertex $v \in V$ such that $T_G(v) = V$ and no other vertex can trace G . And since $\tau(G) = 1$ G must be a track connected floral graph [4]. \square

Remark 2.1. For any graph G ,

$$t(G, 1) = \begin{cases} 1 & \text{if } G \text{ is track connected floral graph} \\ |V| & \text{if } G \text{ is track connected} \\ 0 & \text{otherwise.} \end{cases}$$

Theorem 2.13. Let G be a graph of order n with r trace free vertices. If $T(G; x) = \sum_{i=\tau(G)}^n t(G, i)x^i$ is its cycle tracking polynomial, then $r = n - t(G, n - 1)$.

Proof. Suppose that $A \subset V(G)$ is the set of all trace free vertices. Then by hypothesis, $|A| = r$. For a vertex $v \in V(G)$, the set $V(G) \setminus \{v\}$ is a cycle tracking set of G if and only if $v \in V(G) \setminus A$. Therefore $t(G, n - 1) = |V(G \setminus A)| = n - r$. Hence the theorem. \square

Theorem 2.14. *Let $G(V, E)$ be a graph of order n . Then $t(G, 1) = |\{v \in V(G) : T_G(v) = V(G)\}|$.*

Proof. For every $v \in V(G)$, $\{v\}$ is a cycle tracking set if and only if v traces all vertices. ie; $T_G(v) = V(G)$. \square

3 Cycle tracking polynomial for some graphs

In this section we find the cycle tracking polynomial for some graphs.

Definition 3.1. [?] A firefly graph $F_{s,t,n-2s-2t-1}$ ($s \geq 0, t \geq 0$ and $n - 2s - 2t - 1 \geq 0$) is a graph of order n that consists of s triangles, t pendent paths of length 2 and $n - 2s - 2t - 1$ pendant edges sharing a common vertex.

Theorem 3.1. $\tau(F_{s,t,n-2s-2t-1}) = n - 2s$.

Proof. The graph $F_{s,t,n-2s-2t-1}$ has $n - 2s - 1$ trace free vertices and the common vertex traces all s triangles. So the $n - 2s - 1$ trace free vertices together with the common vertex form a τ -set. Hence $\tau(F_{s,t,n-2s-2t-1}) = n - 2s$. \square

Theorem 3.2. $t(F_{s,t,n-2s-2t-1}, p) = \binom{2s}{p-n+2s}$ if $n - 2s \leq p \leq n - s - 2$

$$\text{and } t(F_{s,t,n-2s-2t-1}, p) = \binom{2s}{p-n+2s} + \sum_{i_1=1}^{p-n+s+2} \binom{2}{i_1} \left[\sum_{i_2=1}^{p-n+s-i_1+3} \binom{2}{i_2} \left[\dots \sum_{i_j=1}^{p-n+s-i_1-i_2-\dots+j+1} \binom{2}{i_j} \left[\dots \left[\sum_{i_{s-1}=1}^{p-n+s-i_1-i_2-\dots-i_{s-1}+s} \binom{2}{i_{s-1}} \left(\binom{2}{p-n+2s+1-i_1-i_2-\dots-i_{s-1}} \right) \dots \right] \dots \right] \right] \right] \text{ if } n - s - 1 \leq p \leq n.$$

Proof. since $\tau(F_{s,t,n-2s-2t-1}, x) = n - 2s$. case(1) if $n - 2s \leq p \leq n - s - 2$ Then any cycle tracking set S contains the common vertex. So the central vertex together with $n - 2s - 1$ trace free vertices and $p - n + 2s$ other vertices constitute a cycle tracking set S and it can be chosen in $\binom{2s}{p-n+2s}$ ways. case(2) if $n - s - 1 \leq p \leq n$

Here the central vertex together with $n - 2s - 1$ trace free vertices and $p - n + 2s$ other constitute a cycle tracking set S and a set of vertices S of cardinality p having at least one element from each triangle is also form a cycle tracking set and it can be chosen in

$$\left(\binom{2s}{p-n+2s} \right)^+ \sum_{i_1=1}^{p-n+s+2} \binom{2}{i_1} \left[\sum_{i_2=1}^{p-n+s-i_1+3} \binom{2}{i_2} \left[\dots \sum_{i_j=1}^{p-n+s-i_1-i_2-\dots+j+1} \binom{2}{i_j} \right. \right. \\ \left. \left. \left[\dots \left[\sum_{i_{s-1}=1}^{p-n+s-i_1-i_2-\dots-i_{s-1}+s} \binom{2}{i_{s-1}} \binom{2}{p-n+2s+1-i_1-i_2-\dots-i_{s-1}} \right] \dots \right] \right] \right] \\ \text{ways.} \quad \square$$

Theorem 3.3. $T(F_{s,t,n-2s-2t-1}, x) = \sum_{p=n-2s}^{n-s-2} \left(\binom{2s}{p-n+2s} \right)^+ \sum_{p=n-s-2}^n \left[\left(\binom{2s}{p-n+2s} \right) \right. \\ \left. + \sum_{i_1=1}^{p-n+s+2} \binom{2}{i_1} \left[\sum_{i_2=1}^{p-n+s-i_1+3} \binom{2}{i_2} \left[\dots \sum_{i_j=1}^{p-n+s-i_1-i_2-\dots+j+1} \binom{2}{i_j} \right. \right. \right. \\ \left. \left. \left[\dots \left[\sum_{i_{s-1}=1}^{p-n+s-i_1-i_2-\dots-i_{s-1}+s} \binom{2}{i_{s-1}} \binom{2}{p-n+2s+1-i_1-i_2-\dots-i_{s-1}} \right] \dots \right] \right] \right] \right] \right]$

Definition 3.2. A Lollipop graph $L_{n,m}$ is obtained by joining K_n to a path P_m of length m with a bridge.

Theorem 3.4. $\tau(L_{n,m}) = m + 1$

Proof. Since a vertex in K_n can trace all vertices in it and all vertices of P_m are trace free vertices we need at least $m + 1$ vertices to trace $L_{n,m}$. Hence $\tau(L_{n,m}) = m + 1$. □

Theorem 3.5. $T(L_{n,m}, x) = ((1 + x)^n - 1)x^m$

Definition 3.3. A Tadpole $T_{(n,l)}$ is a graph obtained by attaching a path P_l to one of the vertices of the cycle C_n by a bridge.

Theorem 3.6. $\tau(T_{(n,l)}) = l + 1$

Proof. Since a vertex in C_n can trace all vertices in it and all vertices of P_l are trace free vertices we need at least $m + 1$ vertices to trace $T_{(n,l)}$ and hence $\tau(T_{(n,l)}) = l + 1$ □

Theorem 3.7. $T(T(n, l)) = ((1 + x)^n - 1)x^l$

Definition 3.4. For a positive integer n, a helm graph, denoted by H_n is obtained from the Wheel W_n by joining a pendant vertex to each vertex in the outer circle of W_n by means of an edge.

Theorem 3.8. $\tau(H_n) = n$

Proof. Since H_n contains $n - 1$ pendant vertices, all these vertices belong to every cycle tracking set. Since W_n is track connected, a vertex of W_n together with the pendant vertices form a minimal cycle tracking set. So $\tau(H_n) = n$. □

Theorem 3.9. $T(H_n) = ((1 + x)^n - 1)x^{n-1}$

Definition 3.5. For a positive integer $n > 3$, a web graph $WB_n, n > 3$ is obtained by joining the pendent vertices of a helm H_n to form a cycle and then adding a single pendent edge to each vertex of this outer cycle. The web graph WB_n has $3n - 2$ vertices and $3(n - 1)$ edges.

Theorem 3.10. $\tau(WB_n) = n$

Theorem 3.11. $T(WB_n) = ((1 + x)^{2n-1} - 1)x^{n-1}$

Definition 3.6. A friendship graph F_n is the one point union of n copies of the cycle C_3 .

Since F_n is a track connected floral graph with n petals each having 3 vertices we have:

Theorem 3.12. $\tau(F_n) = 1$

Using Theorem 2.6:

Theorem 3.13. $t(F_n, i) = \begin{cases} \binom{2n}{i-1} & 1 \leq i \leq n-1 \\ \binom{2n}{i-1} + \binom{n}{i-n}2^{2n-i} & n \leq i \leq 2n \end{cases}$
 and $T(F_n, x) = x + 2nx^2 + \dots + \binom{2n}{i-1}x^i + \dots + \binom{2n}{n-2}x^{n-1}$
 $+ \left[\binom{2n}{n-1} + 2^n \right] x^n + \dots + \left[\binom{2n}{j-1} + \binom{n}{j-n}2^{2n-j} \right] x^j + \dots + x^{2n+1}.$

Definition 3.7. An armed crown $C_n \odot P_m$ is a graph obtained by attaching path P_m to every vertex of the cycle C_n .

Theorem 3.14. $\tau(C_n \odot P_m) = mn + 1$

Proof. Since a vertex in C_n can trace all vertices in it and the remaining mn vertices are trace free vertices. So we need at least $mn + 1$ vertices to trace $C_n \odot P_m$. Hence $\tau(C_n \odot P_m) = mn + 1$ □

Theorem 3.15. $T(C_n \odot P_m) = ((1 + x)^n)x^{mn}$

The corona of two graphs G_1 and G_2 is the graph $G = G_1 \circ G_2$ formed from one copy of G_1 and $|V(G_1)|$ copies of G_2 , where the i^{th} vertex of G_1 is adjacent to every vertex in the i^{th} copy of G_2 [?].

Theorem 3.16. For any graph G of order n , $\tau(G \circ K_1) = n + \tau(G)$.

Theorem 3.17. $T(G \circ K_1, x) = T(G, x)x^n$

Theorem 3.18. Let G be a graph with n vertices and H be a connected graph with $m(m > 1)$ vertices. Then $\tau(G \circ H) = n$.

In particular if $G = K_1$, then $\tau(K_1 \circ H) = 1$

Corollary 3.19. For a connected graph H of cardinality $m(m > 1)$, $T(K_1 \circ H, x) = (1 + x)^{m+1} - 1$.

Proof. Since $K_1 \circ H$ is a track connected graph with $m + 1$ vertices, $T(K_1 \circ H, x) = (1 + x)^{m+1} - 1$ \square

Theorem 3.20. Let G be a graph with n vertices and H be a connected graph with $m(m > 1)$ vertices. Then $T(G \circ H, x) = [(1 + x)^{m+1} - 1]^n$.

Proof. It is enough to prove for $n=2$

For $k \geq \tau(G \circ H) = 2$, a cycle tracking set of k vertices in $G \circ H$ is chosen by selecting j ($1 \leq j \leq k - 1$) vertices from first copy of $K_1 \circ H_m$ and $k - j$ vertices from second copy of $K_1 \circ H_m$. The number of way of doing this over all $k = 2, 3, \dots, mn$ is exactly the coefficient of x^k in $[(1 + x)^{m+1} - 1]^2$. So $T((G \circ H), x) = [(1 + x)^{m+1} - 1]^2$. \square

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